Welcome!

Optimization & Vectorization

J. Bikker - Sep-Nov 2019 - Lecture 4: “Caching (2)"
Today’s Agenda:

- Caching: Recap
- Data Locality
- Alignment
- False Sharing
- A Handy Guide  *(to Pleasing the Cache)*
Recap

Refresher:

Three types of cache:

- Fully associative
- Direct mapped
- N-set associative

In an N-set associative cache, each memory address can be stored in N slots.

Example:
- 32KB, 8-way set-associative, 64 bytes per cache line: 64 sets of 512 bytes.
Recap

**32KB, 8-way set-associative, 64 bytes per cache line:**

*64 sets of 512 bytes*

![32-bit address diagram]

- **tag**: 31 bits
- **set nr**: 12 bits
- **offset**: 5 bits
Recap

32KB, 8-way set-associative, 64 bytes per cache line: 64 sets of 512 bytes

![32-bit address diagram]

Examples:

- **0x00001234**: 0001 001000 110100
- **0x00008234**: 1000 001000 110100
- **0x00006234**: 0110 001000 110100
- **0x0000A234**: 1010 001000 110100
- **0x0000A240**: 1010 001001 000000
- **0x0000F234**: 1111 001000 110100
Recap

32KB, 8-way set-associative, 64 bytes per cache line: 64 sets of 512 bytes

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32-bit address

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Recap

**32KB, 8-way set-associative, 64 bytes per cache line:**

64 sets of 512 bytes

<table>
<thead>
<tr>
<th>31</th>
<th>12</th>
<th>11</th>
<th>6</th>
<th>5</th>
<th>0</th>
</tr>
</thead>
<tbody>
<tr>
<td>tag</td>
<td></td>
<td>set nr</td>
<td></td>
<td>offset</td>
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</tr>
</tbody>
</table>

32-bit address

**Theoretical consequence:**

- Address 0, 4096, 8192, ... map to the same set (which holds max. 8 addresses)
- Consider `int value[1024][1024]:`
  - `value[0, 1, 2...][x]` map to the same set
  - Querying this array vertically:
    - Will quickly result in evictions
    - Will use only 512 bytes of your cache
64 bytes per cache line

Theoretical consequence:

- If address $X$ is pulled into the cache, so is $(X+1,..., X+63)$.

Example*:

```cpp
int arr = new int[64 * 1024 * 1024];
// loop 1
for(int i = 0; i < 64 * 1024 * 1024; i++) arr[i] *= 3;
// loop 2
for(int i = 0; i < 64 * 1024 * 1024; i += 16) arr[i] *= 3;
```

Which one takes longer to execute?

64 bytes per cache line

Theoretical consequence:

- If address \( X \) is removed from cache, so is \( (X+1 \ldots X+63) \).
- If the object you’re querying straddles the cache line boundary, you may suffer not one but \textit{two} cache misses.

Example:

```c
struct Pixel { float r, g, b; }; // 12 bytes
Pixel screen[768][1024];
```

Assuming pixel \((0,0)\) is aligned to a cache line boundary, the offsets in memory of pixels \((0,1..5)\) are 12, 24, 36, 48, 60, \ldots . Walking column 5 will be very expensive.
Recap

Considering the Cache

- Size
- Cache line size and alignment
- Aliasing
- Sharing
- Access patterns
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Why do Caches Work?

1. Because we tend to reuse data.
2. Because we tend to work on a small subset of our data.
3. Because we tend to operate on data in patterns.
Data Locality

- **Very short term**: variable ‘i’ being used intensively in a loop ➔ register
- **Short term**: lookup table for square roots being used on every input element ➔ L1 cache
- **Mid-term**: particles being updated every frame ➔ L2, L3 cache
- **Long term**: sound effect being played ~ once a minute ➔ RAM
- **Very long term**: playing the same CD every night ➔ disk
**Data Locality**

**Reusing data**

**Ideal pattern:**
- load data sequentially.

**Typical pattern:**
- *whatever the algorithm dictates.*
Data Locality

Example: rotozooming
Data Locality

Example: rotozooming
Data Locality

Example: rotozooming

Improving data locality: z-order / Morton curve

Method:

X = 1 1 0 0 0 1 0 1 1 0 1 1 0 1 1 0 1

Y = 1 0 1 1 0 1 1 0 1 0 1 0 1 1 1 0 0

M = 11011010001110011001111101
Data Locality

Wikipedia:

**Temporal Locality** – “If at one point in time a particular memory location is referenced, then it is likely that the same location will be referenced again in the near future.”

**Spatial Locality** – “If a particular memory location is referenced at a particular time, then it is likely that nearby memory locations will be referenced in the near future.”

Data Locality

How do we increase data locality?

**Linear access** – Sometimes as simple as swapping for loops *

**Tiling** – Example of working on a small subset of the data at a time.

**Streaming** – Operate on/with data until done.

**Reducing data size** – Smaller things are closer together.

* For an elaborate example see [https://www.cs.duke.edu/courses/cps104/spring11/lects/19-cache-sw2.pdf](https://www.cs.duke.edu/courses/cps104/spring11/lects/19-cache-sw2.pdf)
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Alignment

Cache line size and data alignment

What is wrong with this struct?

```c
struct Particle {
    float x, y, z;
    float vx, vy, vz;
    float mass;
};
// size: 28 bytes
```

Better:

```c
struct Particle {
    float x, y, z;
    float vx, vy, vz;
    float mass, dummy;
};
// size: 32 bytes
```

Note:

As soon as we read any field from a particle, the other fields are guaranteed to be in L1 cache.

If you update x, y and z in one loop, and vx, vy, vz in a second loop, it is better to merge the two loops.

Two particles will fit in a cache line (taking up 56 bytes).

The next particle will be in two cache lines.
Alignement

Cache line size and data alignment

What is wrong with this allocation?

```c
struct Particle
{
   float x, y, z;
   float vx, vy, vz;
   float mass, dummy;
};
// size: 32 bytes
Particle particles[512];
```

Although two particles will fit in a cache line, we have no guarantee that the address of the first particle is a multiple of 64.

Note:

Is it bad if particles straddle a cache line boundary?

Not necessarily: if we read the array sequentially, we sometimes get 2, but sometimes 0 cache misses.

For random access, this is not a good idea.
Cache line size and data alignment

Controlling the location in memory of arrays:

An address that is dividable by 64 has its lowest 6 bits set to zero. In hex: all addresses ending with 40, 80 and C0.

Enforcing this:

```c
Particle* particles = _aligned_malloc(512 * sizeof(Particle), 64);
```

Or:

```c
__declspec(align(64)) struct Particle { ... };
```
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False Sharing

Multiple Cores using Caches

Two cores can hold copies of the same data.

Not as unlikely as you may think – Example:

```c
byte data = new byte[COUNT];
for (int i = 0; i < COUNT; i++)
    data[i] = rand() % 256;

// count byte values
int counter[256];
for (int i = 0; i < COUNT; i++)
    counter[byteArray[i]]++;
```
Today's Agenda:

- Caching: Recap
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- A Handy Guide *(to Pleasing the Cache)*
Easy Steps

How to Please the Cache

Or: “how to evade RAM”

1. Keep your data in registers

Use fewer variables
Limit the scope of your variables
Pack multiple values in a single variable
Use floats and ints (they use different registers)
Compile for 64-bit (more registers)
Arrays will never go in registers

Unions will never go in registers
Easy Steps

How to Please the Cache

Or: “how to evade RAM”

2. Keep your data local

Read sequentially
Keep data small
Use tiling / Morton order
Fetch data once, work until done (streaming)
Reuse memory locations
How to Please the Cache

Or: “how to evade RAM”

3. Respect cache line boundaries

Use padding if needed
Don’t pad for sequential access
Use aligned malloc / __declspec align
Assume 64-byte cache lines
How to Please the Cache

Or: “how to evade RAM”

4. Advanced tricks

Prefetch
Use a prefetch thread (theoretical...)

Use *streaming writes*

Separate mutable / immutable data
Easy Steps

How to Please the Cache

Or: “how to evade RAM”

5. Be informed

Use the profiler!
Today's Agenda:

- Caching: Recap
- Data Locality
- Alignment
- False Sharing
- A Handy Guide (to Pleasing the Cache)
END of “Caching (2)”

next lecture: “SIMD (1)”
Practical
1. Timing the Rotozoomer
Data Locality

Example: rotozooming
Data Locality

Example: rotozooming
2. Simple Access Pattern
Data Locality

Example: dotcloud
3. LUTs Gone Wrong
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<th>Source</th>
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<th>Instruction</th>
<th>Address</th>
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End of Practical